

CSE 417T: Introduction to Machine Learning

Lecture 5: VC-Dimension

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Recall: Hoeffding's Inequality

- $P\{|E_{in}(g) - E_{out}(g)| > \epsilon\} \leq 2(m)e^{-2\epsilon^2 n}$ where $m = |\mathcal{H}|$
- Growth function: $m_{\mathcal{H}}(n) = \max_{(\vec{x}_1, \dots, \vec{x}_n) \in \mathcal{X}} |\mathcal{H}(\vec{x}_1, \dots, \vec{x}_n)|$
- Can we simply replace m in Hoeffding's inequality with the growth function?

Vapornik- Chervonenkis (VC)-Bound

$$P\{|E_{in}(g) - E_{out}(g)| > \epsilon\} \leq 4m_{\mathcal{H}}(2n)e^{-\frac{1}{8}\epsilon^2 n}$$



$$E_{out}(g) \leq E_{in}(g) + \sqrt{\frac{8}{n} \log\left(\frac{4m_{\mathcal{H}}(2n)}{\delta}\right)}$$

with probability at least $1 - \delta$

Recall

- Claim: if $\exists k$ s.t. k is a breakpoint for \mathcal{H} , then $m_{\mathcal{H}}(n)$ is bounded by a polynomial in n
- If $m_{\mathcal{H}}(n)$ is bounded by a polynomial in n , then

$$\sqrt{\frac{8}{n} \log \left(\frac{4m_{\mathcal{H}}(2n)}{\delta} \right)} \rightarrow 0 \text{ as } n \rightarrow \infty$$

Bounding $m_{\mathcal{H}}(n)$

- Let $B(n, k) =$ the maximum number of dichotomies on n points s.t. no subset of k points is shattered
- If k is a breakpoint for \mathcal{H} , then $m_{\mathcal{H}}(n) \leq B(n, k)$
- If $B(n, k)$ is bounded by a polynomial in n and $m_{\mathcal{H}}(n)$ is bounded by $B(n, k)$, then $m_{\mathcal{H}}(n)$ is bounded by a polynomial in n

$B(n, k)$: Example

- What is $B(4, 2)$?

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
+1	+1	+1	+1
+1	+1	+1	-1
+1	+1	-1	+1
+1	-1	+1	+1
-1	+1	+1	+1

$B(n, k)$: Example

- What is $B(4, 2)$?

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
+1	+1	+1	+1
+1	+1	+1	-1
+1	+1	-1	+1
+1	-1	+1	+1
-1	+1	+1	+1
-1	-1	+1	+1

$B(n, k)$: Example

- What is $B(4, 2)$?

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
+1	+1	+1	+1
+1	+1	+1	-1
+1	+1	-1	+1
+1	-1	+1	+1
-1	+1	+1	+1
+1	-1	-1	-1

$B(n, k)$:
Example

- $B(4, 2) = 5$

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
+1	+1	+1	+1
+1	+1	+1	-1
+1	+1	-1	+1
+1	-1	+1	+1
-1	+1	+1	+1

$B(n, k)$:
Example

- $B(4, 2) = 5$

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
+1	+1	+1	+1
+1	-1	-1	-1
+1	+1	-1	+1
+1	-1	-1	+1
-1	+1	+1	+1

$B(n, k)$:
Example

- $B(3, 2) = 4$

\vec{x}_1	\vec{x}_2	\vec{x}_3
+1	+1	+1
+1	+1	-1
+1	-1	+1
-1	+1	+1

$B(n, k)$: Example

- $B(4, 3) = 11$

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
-1	-1	+1	+1
-1	+1	-1	+1
+1	-1	-1	-1
+1	+1	+1	+1
-1	+1	+1	+1
+1	-1	+1	+1
+1	+1	-1	+1
+1	+1	+1	-1
-1	+1	+1	-1
+1	-1	+1	-1
+1	+1	-1	-1

$B(n, k)$: Example

- $B(4, 3) = \alpha + 2\beta$

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
-1	-1	+1	+1
-1	+1	-1	+1
+1	-1	-1	-1
+1	+1	+1	+1
-1	+1	+1	+1
+1	-1	+1	+1
+1	+1	-1	+1
+1	+1	+1	-1
-1	+1	+1	-1
+1	-1	+1	-1
+1	+1	-1	-1

$S_1, |S_1| = \alpha$

$S_2^+, |S_2^+| = \beta$

$S_2^-, |S_2^-| = \beta$

$B(n, k)$: Example

- $B(4, 3) = \alpha + 2\beta$

- $\alpha + \beta \leq B(3, 3)$

- $\beta \leq B(3, 2)$

- $B(4, 3) \leq B(3, 3) + B(3, 2)$

\vec{x}_1	\vec{x}_2	\vec{x}_3	\vec{x}_4
-1	-1	+1	+1
-1	+1	-1	+1
+1	-1	-1	-1
+1	+1	+1	+1
-1	+1	+1	+1
+1	-1	+1	+1
+1	+1	-1	+1
+1	+1	+1	-1
-1	+1	+1	-1
+1	-1	+1	-1
+1	+1	-1	-1

$S_1, |S_1| = \alpha$

$S_2^+, |S_2^+| = \beta$

$S_2^-, |S_2^-| = \beta$

$B(n, k)$: Recursion

- $B(n, k) \leq B(n - 1, k) + B(n - 1, k - 1)$

- $B(n, k) = \alpha + 2\beta$

\vec{x}_1	\vec{x}_2	...	\vec{x}_{n-1}	\vec{x}_n	
-1	-1	...	+1	-1	} $S_1, S_1 = \alpha$
-1	+1	...	-1	+1	
⋮	⋮	⋮	⋮	⋮	
+1	-1	...	-1	-1	} $S_2^+, S_2^+ = \beta$
+1	+1	...	+1	+1	
-1	+1	...	+1	+1	
⋮	⋮	⋮	⋮	⋮	
+1	-1	...	+1	+1	} $S_2^-, S_2^- = \beta$
+1	+1	...	+1	-1	
-1	+1	...	+1	-1	
⋮	⋮	⋮	⋮	⋮	
+1	-1	...	+1	-1	

$B(n, k)$: Recursion

- $B(n, k) \leq B(n - 1, k) + B(n - 1, k - 1)$

- $B(n, k) = \alpha + 2\beta$

- $\alpha + \beta \leq B(n - 1, k)$

\vec{x}_1	\vec{x}_2	...	\vec{x}_{n-1}	\vec{x}_n	
-1	-1	...	+1	-1	} $S_1, S_1 = \alpha$
-1	+1	...	-1	+1	
⋮	⋮	⋮	⋮	⋮	
+1	-1	...	-1	-1	} $S_2^+, S_2^+ = \beta$
+1	+1	...	+1	+1	
-1	+1	...	+1	+1	
⋮	⋮	⋮	⋮	⋮	
+1	-1	...	+1	+1	
+1	+1	...	+1	-1	
-1	+1	...	+1	-1	
⋮	⋮	⋮	⋮	⋮	
+1	-1	...	+1	-1	

$B(n, k)$: Recursion

- $B(n, k) \leq B(n - 1, k) + B(n - 1, k - 1)$
- $B(n, k) = \alpha + 2\beta$
- $\alpha + \beta \leq B(n - 1, k)$
- $\beta \leq B(n - 1, k - 1)$

\vec{x}_1	\vec{x}_2	...	\vec{x}_{n-1}	\vec{x}_n
-1	-1	...	+1	-1
-1	+1	...	-1	+1
⋮	⋮	⋮	⋮	⋮
+1	-1	...	-1	-1
+1	+1	...	+1	+1
-1	+1	...	+1	+1
⋮	⋮	⋮	⋮	⋮
+1	-1	...	+1	+1
+1	+1	...	+1	-1
-1	+1	...	+1	-1
⋮	⋮	⋮	⋮	⋮
+1	-1	...	+1	-1

} $S_2^-, |S_2^-| = \beta$

Bounding $B(n, k)$

- Claim: $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$

- Intuition: if recursion, then induction

Bounding $B(n, k)$

- Claim: $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$
- Induction (base cases):
 - $\forall n \geq 1, B(n, 1) = 1 \leq \sum_{i=0}^{k-1} \binom{n}{i} = \binom{n}{0} = 1$
 - $\forall k \geq 2, B(1, k) = 2 \leq \sum_{i=0}^{k-1} \binom{1}{i} = \binom{1}{0} + \binom{1}{1} = 2$

Bounding $B(n, k)$

- Claim: $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$
- Induction (inductive step):
- Suppose $B(n', k') \leq \sum_{i=0}^{k'-1} \binom{n'}{i} \quad \forall n' < n, k' < k$
- Prove $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$

Bounding $B(n, k)$

- Prove $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$
- $B(n, k) \leq B(n-1, k) + B(n-1, k-1)$
$$\leq \sum_{i=0}^{k-1} \binom{n-1}{i} + \sum_{i=0}^{k-2} \binom{n-1}{i}$$
$$= 1 + \sum_{i=1}^{k-1} \binom{n-1}{i} + \sum_{i=1}^{k-1} \binom{n-1}{i-1}$$
$$= 1 + \sum_{i=1}^{k-1} \left(\binom{n-1}{i} + \binom{n-1}{i-1} \right) = 1 + \sum_{i=1}^{k-1} \binom{n}{i}$$

Bounding $B(n, k)$

- Prove $B(n, k) \leq \sum_{i=0}^{k-1} \binom{n}{i}$
- $B(n, k) \leq B(n-1, k) + B(n-1, k-1)$
$$\leq \sum_{i=0}^{k-1} \binom{n-1}{i} + \sum_{i=0}^{k-2} \binom{n-1}{i}$$
$$= 1 + \sum_{i=1}^{k-1} \binom{n-1}{i} + \sum_{i=1}^{k-1} \binom{n-1}{i-1}$$
$$= 1 + \sum_{i=1}^{k-1} \left(\binom{n-1}{i} + \binom{n-1}{i-1} \right) = \sum_{i=0}^{k-1} \binom{n}{i} \quad \blacksquare$$